

# Social Choice Theory

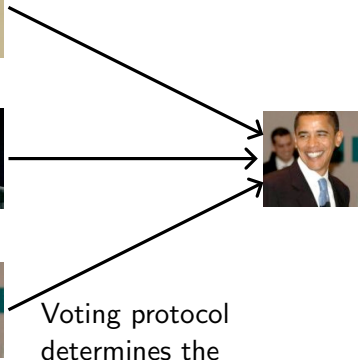
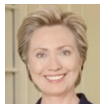
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# Introduction

- Social choice: preference aggregation
- Our settings
  - ▶ A set of agents have preferences over a set of alternatives
  - ▶ Taking preferences of all agents, the mechanism outputs a social preference over the set of alternatives or output a single winner
  - ▶ Hope to satisfy some desired properties
- Voting protocols are examples of social choice mechanisms
- Readings: SLB 9.1 – 9.4

# Voting



Voting protocol determines the winner or the final ordering

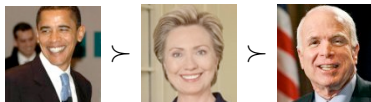
# Example Voting Protocols

- Plurality Voting
  - ▶ Each voter cast a single vote.
  - ▶ The candidate with the most votes is selected.
- Approval Voting
  - ▶ Each voter can cast a single vote for as many candidates as he wants.
  - ▶ The candidate with the most votes is selected.
- Single Transferable Vote (Instant Roundoff)
  - ▶ Each candidate votes for their most-preferred candidate
  - ▶ The candidate with the fewest votes is eliminated
  - ▶ Each voter who voted for the eliminated candidate transfers their vote to their most-preferred candidate among the remaining candidates
- Borda Voting
  - ▶ Each voter submits a full ordering on the  $m$  candidates
  - ▶ Candidates of an ordering get score  $(m - 1, m - 2, \dots, 0)$
  - ▶ The candidate with the highest score is selected

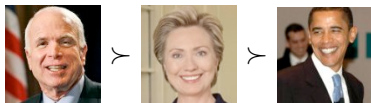
# Pairwise Elections



2 prefer Obama to McCain



2 prefer McCain to Hillary



2 prefer Obama to Hillary



# More Voting Protocols

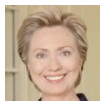
- Pairwise elimination
  - ▶ Pair candidates with a schedule
  - ▶ The candidate who is preferred by a minority of voters is deleted
  - ▶ Repeat until only one candidate is left
- Slater
  - ▶ The overall ordering that is inconsistent with as few pairwise elections as possible is selected.
  - ▶ NP-hard
- Kemeney
  - ▶ The overall ordering that is inconsistent with as few votes on pairs of candidates as possible.
  - ▶ NP-hard
- ... and many other voting rules

What is the perfect voting protocol?

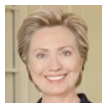
# Condorcet Condition

- A candidate is a **Condorcet winner** if it wins all its pairwise elections.
- A voting protocol satisfies the **Condorcet condition**, if the Condorcet winner, if exists, must be elected by the protocol.
- Condorcet winner may not exist.
- Many voting protocols do not satisfy the Condorcet condition.

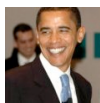
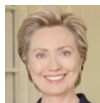
# Condorcet Circle



2 prefer Obama to McCain



2 prefer McCain to Hillary



2 prefer Hillary to Obama

?



# An Example of Condorcet Condition

499 agents:  $a \succ b \succ c$

3 agents:  $b \succ c \succ a$

498 agents:  $c \succ b \succ a$

- Which candidate is the Condorcet winner if exists?
- Which candidate is the plurality voting selected?
- Which candidate is the Single Transferable Vote selected?

# Voting Paradox: Sensitivity to A Losing Candidate

35 agents:  $a \succ c \succ b$

33 agents:  $b \succ a \succ c$

32 agents:  $c \succ b \succ a$

- Which alternative is the winner under plurality voting?
- Which alternative is the winner under Borda voting?
- What happens if  $c$  drops off?

# Notations

- $N$ : a set of individuals,  $|N| = n$
- $A$ : a set of alternatives,  $|A| = m$
- $\succ_i$ : agent  $i$ 's preference over  $A$  (e.g.  $a_i \succ_i a_3 \succ_i a_5$ )
- $L$ : the set of total orders,  $\succ \in L$
- $L^n$ : the set of preference profiles,  $[\succ] \in L^n$
- A **social welfare function** is a function  $W : L^n \rightarrow L$
- $\succ_W$ : the preference ordering selected by  $W$
- A **social choice function** is a function  $C : L^n \rightarrow A$

# Social Welfare Function: Pareto Efficiency

- A social welfare function  $W$  is **Pareto efficient** if for any  $a_1, a_2 \in A$ ,  $\forall a_1 \succ_i a_2$  implies that  $a_1 \succ_W a_2$ .
- It means that when all agents agree on the ordering of two alternatives, the social welfare function must select the ordering.

# Social Welfare Function: Independence of Irrelevant Alternatives (IIA)

- A social welfare function  $W$  is **independent of irrelevant alternatives** if, for any  $a_1, a_2 \in A$  and any two preference profiles  $[\succ']$ ,  $[\succ''] \in L^n$ ,  $\forall i$

$$(a_1 \succ'_i a_2 \text{ if and only if } a_1 \succ''_i a_2) \Rightarrow \\ (a_1 \succ_{W([\succ'])} a_2 \text{ if and only if } a_1 \succ_{W([\succ''])} a_2).$$

- IIA means that if (1)  $W$  ranks  $a_1$  ahead of  $a_2$  now, and (2) we change the preferences without change the relative preferences between  $a_1$  and  $a_2$ , then  $a_1$  is still ranked ahead of  $a_2$ .
- An example with plurality voting protocol

$$\begin{array}{ll} 499 \text{ agents: } a \succ b \succ c & a \succ b \succ c \\ 3 \text{ agents: } b \succ c \succ a & \Rightarrow b \succ c \succ a \\ 498 \text{ agents: } c \succ b \succ a & b \succ a \succ c \end{array}$$

- None of our rules satisfy IIA

# Social Welfare Function: Nondictatorship

- We do not have a **dictator** if there does not exist an  $i$  such that  $\forall a_1, a_2,$

$$a_1 \succ_i a_2 \Rightarrow a_1 \succ_W a_2$$

- Nondictatorship means that there does not exist a voter such that the social welfare function  $W$  always output the voter's preference

# Arrow's Impossibility Results (1951)

- If  $|A| \geq 3$ , any social welfare function  $W$  can not simultaneously satisfy
  - ▶ Pareto efficiency
  - ▶ Independence of irrelevant alternatives
  - ▶ Nondictatorship
- Most influential result in social choice theory
- Read the proof

Maybe asking for a complete ordering is too much? Let's consider social choice functions.

# Social Choice Function: Weak Pareto Efficiency

- A social choice function  $C$  is **weakly Pareto efficient** if for any preference profile  $[\succ] \in L^n$ , if there exist a pair of alternatives  $a_1$  and  $a_2$  such that  $\forall i \in N, a_1 \succ_i a_2$ , then  $C([\succ]) \neq a_2$ .
- It means that a dominated alternative can not be selected.
- Weak Pareto efficiency implies **unanimity**: If  $a_1$  is the top choice for all agents, we must have  $C([\succ]) = a_1$ .
- Pareto efficient rules satisfy weak Pareto efficiency. But the reverse is not true.



# Social Choice Function: Strong Monotonicity

- A social choice function  $C$  is **strongly monotonic**, if for any preference profile  $[\succ]$  with  $C[\succ] = a$ , then for any other preference profile  $[\succ']$  with the property that

$$\forall i \in N, \forall a' \in A, a \succ_i' a' \text{ if } a \succ_i a',$$

it must be that  $C[\succ'] = a$ .

- Strong monotonicity means that if
  - The current winner is  $a$
  - We change the preference profile in the way such that for if alternative  $a'$  ranks below  $a$  previously it is still below  $a$  in the new preferenceThen,  $a$  is the winner for the new preference profile.

- An example with STV

9 agents:  $a \succ b \succ c$

9 agents:  $b \succ c \succ a$

7 agents:  $c \succ a \succ b$

12 agents:  $a \succ b \succ c$

$\Rightarrow$  6 agents:  $b \succ c \succ a$

7 agents:  $c \succ a \succ b$

- None of our rules satisfy strong monotonicity

## Social Choice Function: Nondictatorship

- A social choice function  $C$  is **nondictatorial** if there does not exist an agent  $i$  such that  $C$  always outputs the top choice of  $i$ .

# Muller-Satterthwaite's Impossibility Results (1977)

- If  $|A| \geq 3$ , any social choice function  $C$  can not simultaneously satisfy
  - ▶ Weak Pareto efficiency (unanimity)
  - ▶ Strong monotonicity
  - ▶ Nondictatorship
- Social choice functions are no simpler than social welfare functions
- Intuition: We can repeatedly probe a social choice function for given pairs of alternatives, and then construct a full social welfare ordering.

# Social Choice Function: Manipulability

- A social choice function is **manipulable** if some voter can be better off by lying about his preference
- An example with plurality voting

1 agent:  $a \succ b \succ c$

2 agents:  $b \succ c \succ a$

2 agents:  $c \succ b \succ a$

## Social Choice Function: Onto

- A social choice function  $C$  is **onto** if for each  $a \in A$  there is a preference profile  $[\succ] \in L^n$  such that  $C([\succ]) = a$ .
- Onto means that every alternative can be a winner under some preference profile.

# Gibbard-Satterthwaite's Impossibility Results (1973, 1975)

- If  $|A| \geq 3$ , any social choice function can not simultaneously satisfy
  - ▶ Nonmanipulable
  - ▶ Onto
  - ▶ Nondictatorship

What's possible?

## Some Possibility Results: Single-Peaked Preferences

- Alternatives are a linear order (e.g. ordered on real line)
- Single-peaked preference: every voter has his most-preferred alternative and prefers alternatives that are closer to his favorite alternative
- Ask the voters to only report his favorite alternative
- The social choice function chooses the **median voter**'s favorite alternative as the winner
- The winner is a Condorcet winner
- Nonmanipulable!